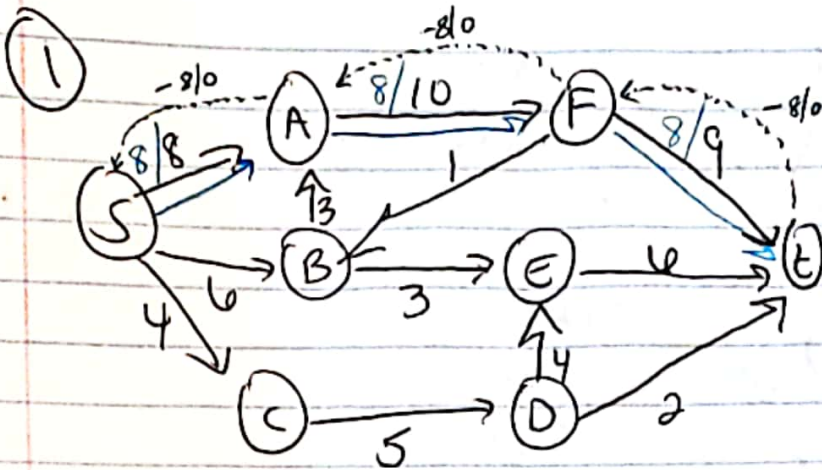
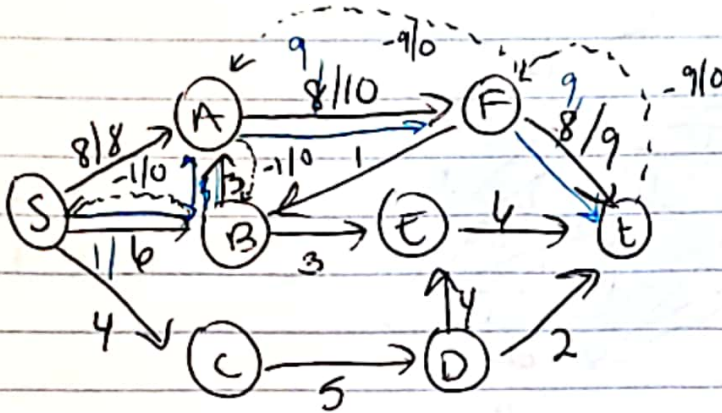


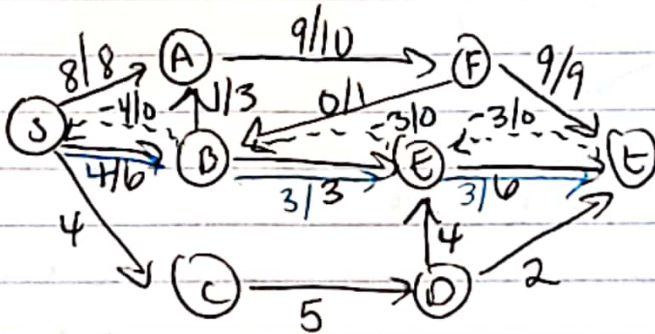
CS 581 Homework #6 solutions



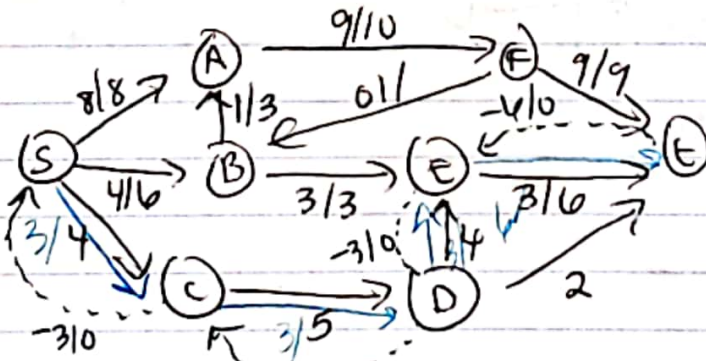
SAFT



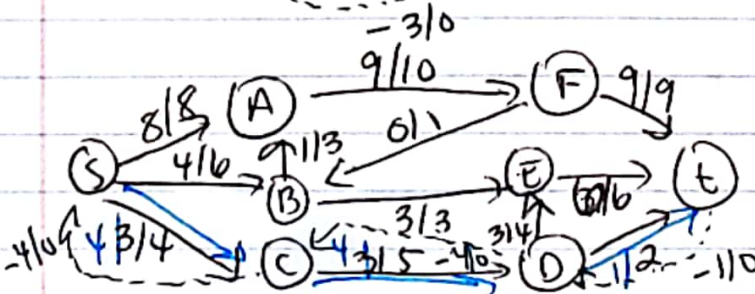
SAFT
SBAFT



SAFT
SBAFT
SBET

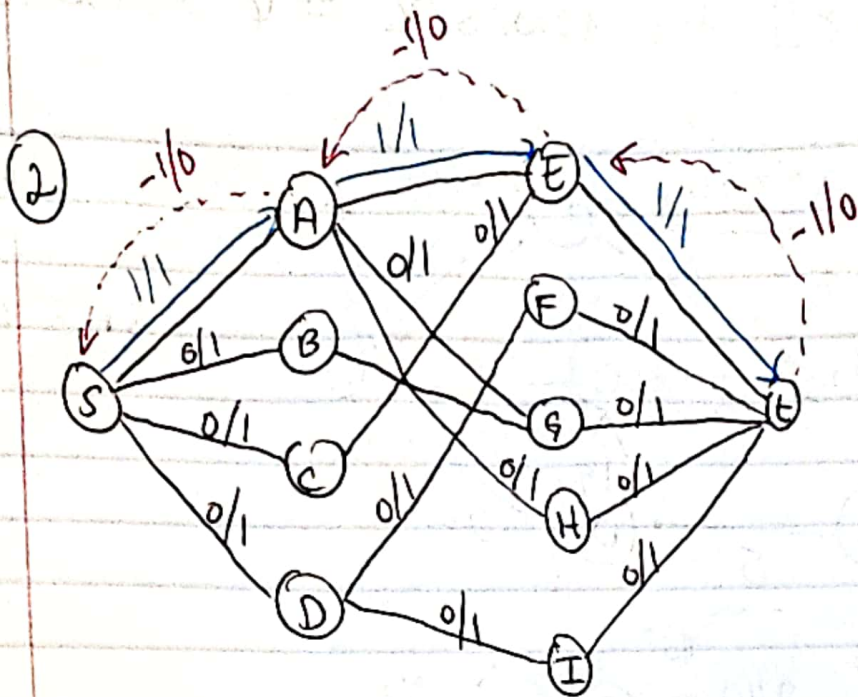


SAFT
SBAFT
SBET
SCDET

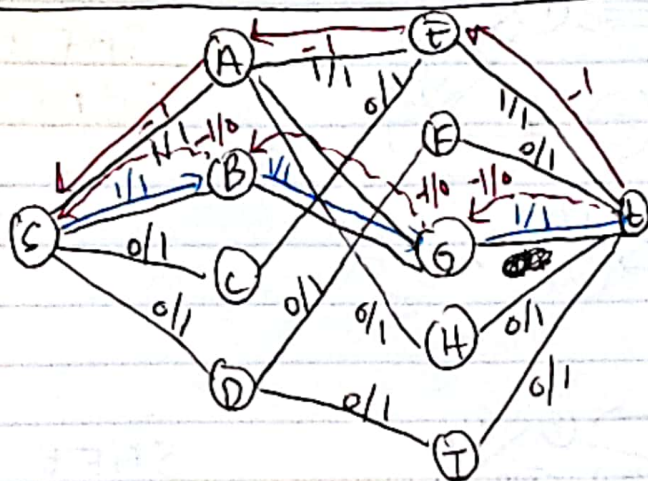


SAFT
SBAFT
SBET
SCDET
SCDET

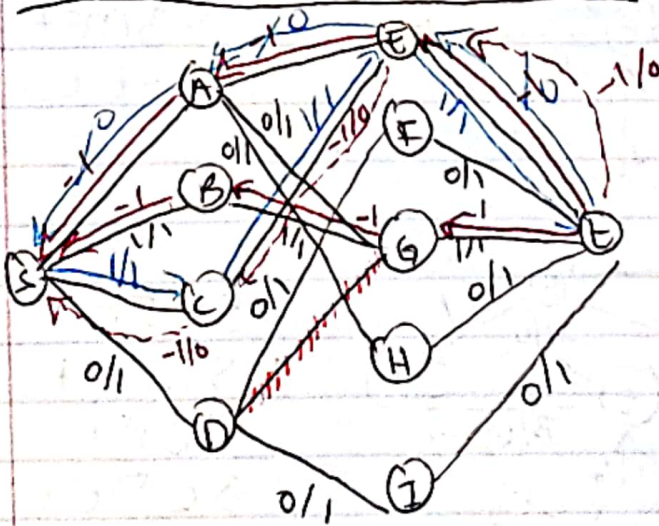
max flow = 16



SAEt

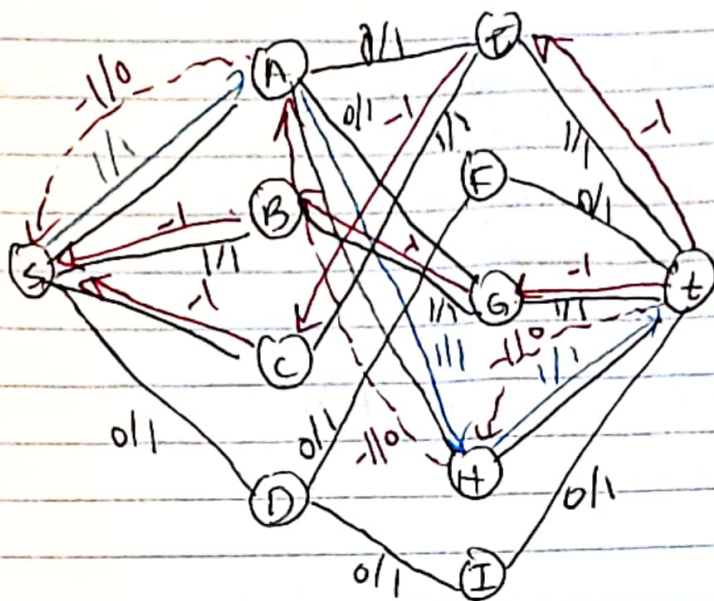


SAEt
SBGt

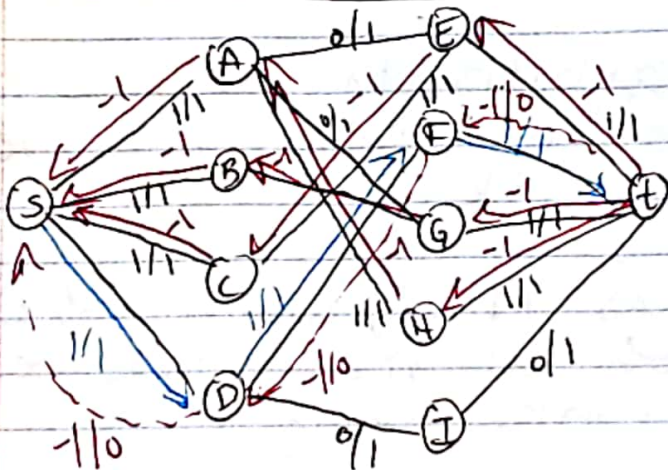


~~SAEt~~
SBGt
SCEt

(next page)



SBGt
 SC Et
 SAHt



SBGt
 SC Et
 SAHt
 SDFt

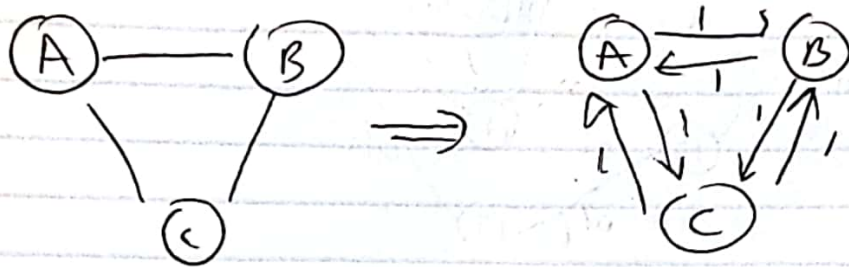
Algorithm terminates!

Final Bipartite matching:

- B → G
- C → E
- A → H
- D → F

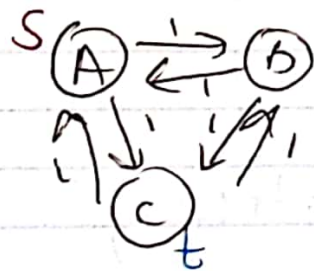
(note I is alternate solution for F)

③ a) We first convert our undirected graph of $O(|E|)$ edges to a directed graph w/ 2 arcs w/ edge capacity 1.



We then arbitrarily fix the source, and allow the other $|V| - 1$ vertices to be the sink.

Perform max flow, where the minimum of the max flows represents the edge connectivity.



b) proof:

Let $G = (V, E)$ be an arbitrary undirected graph.

Let $G' = (V, E')$ be G 's bidirectional equivalent s.t.

$$E' = \{c(u, v), c(v, u) = 1 \quad \forall (u, v) \in E\}$$

Let S be an arbitrary $v \in V$.

Let T be $(V - \{S\})$.

Let C be a set of k edges \Rightarrow removing k disconnects G' .

Hence, $G' - C$ is the disconnected graph (S, T) where $S \in S$.

Then, for each $t \in T$ in G' find the max flow from S to t .

Hence, for at least one $t \in T$ in G' , $t \notin S$, $t \in T$, and $t = k$. (by min cut)

□